Implementation of Load Balancing Algorithm with Cloud Collaboration for Logistics

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Abstract: Cloud computing provides a facility of any data center at any location in the world. There are various applications used in a centralized and distributed data center that suppliers and users can purchase, sell and rent for their information and product accessing. The suppliers and user have no knowledge about where these data centers are located and how they will be operated or maintained by using cloud collaboration. They only know how to connect or use the applications of cloud to perform their jobs. Today, there is lot of demand of cloud in various companies to establish their own data center. In logistics management, sharing of information at different level by different logistics partners, users and suppliers is a big challenge. They all want that every type of information should be shared in real time without any delay at minimum cost. For this, they need cloud based load balancing approaches to control network traffic and overloading on the data center. Load balancing is a distributed technique of workload for load balancing between two or more cloud servers. Load balancing have always task to optimize resource use, virtual machine maintains data center cost, maximize, throughput, minimize response time and minimize overload. There are so many load balancing algorithms.

Key words: Cloud computing, logistics information system, load balancing algorithms, proposed DSBP algorithm, utilized, applications

INTRODUCTION

Cloud computing is exploring dynamic and scalable computing and storage platform which creates needs of users because of its easily adoption. Cloud plays an important role in public sector organization for reducing cost for using information technologies and services (Alsanea et al., 2014). Centralized data center is repository for information distribution and information storage. The centralized data center provides highly efficient for transferring of information and its services for managing all levels of logistics information over the cloud. Logistics management faces different types of issues such as accessing of information at different levels, logistics operations processing, distribution of information inventory and changing in users demand (Feng, 2013). Distribution of any type of information among the suppliers, partners and users is very big issue because number of users is increasing to use various websites over the internet which creates overloading and traffic for providing efficient information for availability of different types of products from different data centers. Load balancing is a distributed technique of workload for balancing between two or more cloud servers. Load balancing considers managing service provided high traffic and always be ready to reduce and balance peak load on the server. Load balancing have always task to optimize resource use, virtual machine maintenance, data center cost, maximum throughput, minimum response time and overload. Load balancing applications redundant mirrored databases in clusters techniques across multiple availability reduces the chance of outages have simultaneously affect the cloud services. If an outage affects one system, the load balancer switches to another available resource. Load balancing technique reduces cost with document management systems and increases availability of resources to reduce amount of downtime that affects business during outages. Load balancing is a new technique which provides facilities of network and resources with maximum throughput and minmum response time. Load balancing provides different solutions applied redundant servers which helps a better distribution of the communication traffic, so that, website availability is conclusively maintained in our daily life.

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Today, load balancing technique is used to enhance the software embedded on compatible hardware in distributed network. Load balancing ensures the implementation of distributed system with restructuring the system workload at highest users and maximum throughput. Load balancing is applied to avoid large difference in workload on servers. There are so many load balancing algorithms utilized to balance the load on servers. In this study, we have presented different load balancing algorithm and evaluated them on the basis of overall response time, data center processing time and total cost of virtual machine and data transfer cost. There are various load balancing algorithms but all existing algorithms have issues to reduce response time, overloading and data transfer cost at same time. We develop a cloud based application by using our proposed distributed service broker policy algorithm for logistics to avoid delay in response time and cost due to multiple users request at same time from different data centers. We present evaluation of some existing load balancing algorithms, round robin algorithm, active monitoring algorithm and throttled load balancing algorithm by using cloud analyst tool and we will also compare the result performance of our propose solution DSBP with these existing algorithms.

Literature review: Load balancing ensures the better performance of distributed system by using redesigning the system overload with highest customer's satisfaction at lower cost. Load balancing techniques are applied for avoiding the situation of different types of over loading on the server. Safiriyu et al. defined Distributed Web Servers (DWS) which provides an effective solution to improve the quality of web services. The experimental result has shown the better performance of network traffic, system throughput, mean response time and system utilization of Cooperative Adaptive Symmetrical Initiated Dynamic/diffusion (CASID) and compared it with PLB (Platform for Load Balancing). Randles et al. (2010) presented a comparison of static and dynamic load balancing algorithms for cloud computing. Li et al. (2010) proposed resource allocation mechanism with preemptable task execution which enhance the utilization of clouds. Singh et al. (2010) proposed Modified Round Robin (MRR) which is better than RR and have less response time and reduced the overhead and saving of memory space (Yaashuwanth and Ramesh, 2010). Yaashuwanth and Ramesh (2010) proposed RR (MRR) algorithm which removes the limitations of simple RR (Singh et al., 2010). Padhy and Rao (2011) defined some existing load balancing algorithms which can be applied over the cloud. Mohanty et al. (2011) also presented dynamic time quantum which modifies with every round of exe cution and result shows that PBDRR performs better than MRR algorithm. Sharma et al. (2012)

Machine (CLBVM) to balance the load in distributed environment but it does not consider fault tolerant system. James and Verma (2012) also analyzed virtual machine load balancing and proposed new virtual machine load balancing algorithm for IaaS framework, implemented of weighted active monitoring load balancing algorithm to achieve better performance in response time and data processing time by using cloudSim. Limbani and Oza (2012) proposed the extended service proximity based routing policy algorithm forrouting of user requests such that cost effective data center selection by using cloud analyst. Mohapatra et al. (2013) analyzed various policies utilization of different algorithms for load balancing by using cloud analyst tool. Elina et al. evaluated particle swarm optimization based cloud scheduler. The study discussed the number of cloud users are able to serve successfully and total number of created Virtual Machines (VMs) in scheduling scenarios, although, the number of intra cloud network messages sent is also evaluated. The simulated results shows some problems related to scheduler succeed to balance the random assignment and genetic algorithms. Bhargava et al. (2013) also analyzed and discussed round robin algorithm used by different data center and calculated the overall response time in better performance. Priya and Subramani (2013) proposed algorithm by using active monitoring load balancing algorithm and resource aware scheduling algorithm for improving resource utilization and scheduled load balancing for high performance in cloud systems. The experiment result of proposed algorithm is the efficient virtual machine is selected for process and minimum execution time of task, it increases the performance and reduces the response time and cost (Priya and Subramani, 2013). Singh and Gangwar (2014) presented comparative study of load balancing algorithms based on the parameters like, response time and cost which are reduced by round robin and active monitoring load balancing algorithm. Mishra and Bhukya (2014) proposed priority and extended priority based round robin service broker algorithms which distribute the requests based on the rating of data centers and gives better performance. Kapgate (2014) also compared the new service broker (DC selection) algorithm with existing service broker algorithm. The proposed algorithm reduced service response time and improved performance of daa center. Nayak and Patel (2015) presented comparison of some existing throttled algorithms and proposed throttled algorithm of load balancing in cloud computing. Both are tested and compared in terms of response time, data center service request time and cost by using cloud analyst. Panwar and Mallick (2015) discussed the various load balancing algorithms and compared them based on parameters like, data processing time and response time, etc. by using

proposed a Central Load Balancing policy for Virtual

round robin and throttled scheduling algorithms in cloud analyst. Bodake (2016) presented a comparison of assorted policies which is used for load reconciliation by using throttled, round robin and First Come First Serve (FCFS) and load aqualization algorithms. This comparison shows that response time was effectively reduced but they did not consider the cost. Khanchi and Tyagi (2016) also proposed and implemented a hybrid approach for virtual machine level load balancing. This algorithm distribute workload among the virtual machines that are available in data center at the same time to minimize the overall response time and data center processing time (Khanchi and Tyagi, 2016). Norwal and Dhingra (2017) proposed a multi objective task scheduling algorithm which consider a wide variety of attributes in cloud environment and uses non dominate sorting for prioritizing the task. The proposed algorithm improved the processing time, cost and average waiting time in comparison with First Come First Serve (FCFS) algorithm (Norwal and Dhingra, 2017). Kanmani and Jayabalan (2017) presented a survey and comparison of existing innovative algorithms based on different parameters like, throughput maitainces, reliability, efficiency, scalability and performance. In this study, we propose a Distributed Service Broker Policy algorithm (DSBP) with the best possible response time, delay and minimum cost in selecting the most suitable data center. DSBP is the mainly implementation of throttled algorithm for taking bandwdth, latency and size of job which achieves minimum response time and minimum cost.

MATERIALS AND METHODS

Issues in existing algorithms: In distributed data center, the main propose of service brokers is to direct the user request to the best data center with better performance because service broker policy has to efficiently select the best data center for the job considering many issues like response time, service time and cost. For directing the user request to the best data center there are so many load balancing policy algorithms like network latency based, service proximity based routing, etc. Sharma et al. (2012) implemented the throttled load balancing algorithm and proposed virtual machine load balancing algorithm which reduce response time effectively but not reduce cost. Bhargava et al. (2013) analyzed round robin algorithm by using different data center and User Base (UB). This result shows that response time is to be minimize but cost should need to be reduce also. Priya and Subramani (2013) proposed a new load balancing algorithm for virtual machine in which virtual machine is selected for process and minimum execution time of task. It reduces the response time and total cost but there is some issues to improve the response time

efficiency parallel to the cost performance. Singh and Gangwar (2014) analyzed the comparison of three existing load balancing algorithm, round robin, active monitoring, throttled algorithm. The experimental result of theses algorithms shows the reduction in response but increment in cost. Mishra and Bhukya (2014) also proposed priority and extended priority based round robin service broker algorithms which distribute the requests based on the rating of data centers and gives better performance in response time but cost is not reduces similar to response time. Kapgate (2014) proposed and compared the new service broker (DC selection) algorithm with existing service broker algorithm. The proposed algorithm reduced service response time but cost is not considered. Nayak and Patel (2015) proposed throttled algorithm and compared it with existing throttled algorithm. The proposed algorithm reduces response time and cost in better performance butif data center will be increased then response time cost will be increased. Panwar and Mallick (2015) also discussed the various load balancing algorithms and compared them based on parameters like, data processing time and response time, etc., by using round robin and throttled scheduling algorithms in cloud analyst. The result of both algorithms shows reduction in response time and cost but if data centers will be increased then response time and cost will increase by using separately round robin and throttled algorithm. Bodake (2016) also presented round Robin algorithm. In this algorithm, the time slits into multiple slices and every node is given a specific time slice. There are multiple requests in multiple processes. Each process is given time slot. If user request completes among time then user must not wait otherwise user need to watch for its next time slot. This may create slow process and overloading. Khanchi and Tyagi (2016) proposed virtual machine load balancing algorithm which is a combination of round robin, throttled, (ESCE) Equally Spread Current Execution and hybrid algorithm. The result shows minimization of overall response time but overall performance is not much improved by using ESCE algorithm. Mehar and Acharya (2017) presented a comprehensive summary of load balancing algorithm in cloud computing. In this study, existing approaches are analyzed to provide load balancing in cloud environment. The comparison of different algorithms presented minimum response time, throughput, fault tolerance, overhead and speed (Mehar and Acharya, 2017). Reema and Sehgal evaluated optimize response time and reconfigure dynamically policies for evaluating task mapping algorithm to minimize resource cost, VM cost, data center cost and processing time but results are not showed successfully. In Table 1, we can see, so many results of different algorithms in perspective of response time and cost to identify the improvement of round robin, active monitoring and throttled algorithms.

	Algorithms			Results	
Researcher's (Years)	Round Robin algorithm	Active monitoring algorithm	Throttledload balancing algorithm	Average response time (m sec)	Total cost (S
Kishore and Thappar (2014)	No	No	No	481.54	Ni1
Behal and Kumar (2014)	Yes	No	No	322.55	2505.48
AliNaser	No	No	Yes	139.98	59.50
Nandwani <i>et al.</i> (2015)	Yes	Yes	Yes	92.10	Ni1
Patel and Rajawat (2015)	Yes	No	Yes	177.59	2.28
Singh et al. (2016a, b)	Yes	No	Yes	57.68	Ni1
Er. Pooja	No	No	No	57.08	12.85
Singh et al. (2016a, b)	Yes	No	No	484.25	Ni1
Shukla and Arora (2016)	Yes	No	Yes	150.87	Ni1
Sharma et al. (2012)	No	No	No	333.09	Ni1
Shipra Gayal	Yes	No	Yes	338.27	Ni1
Manasrahet al. (2017)	Yes	Yes	Yes	498.93	973.53

Table 1: Existing algorithm simulation results (Kishore and Thapper, 2014; Behal and Kumar, 2014; Nandwani et al., 2015; Patel and Rajawat, 2013; Singh et al., 2016a, b; Shukla and Arora, 2016; Sharma and Gupta, 2016

Table 2: Methods of implemented load balancing algorithms

Algori thms	Round Robin	Active Monitoring	Throttled load balancing	DSBP Proposed Algorithm
Metho ds	<pre>public int getNextAyail ableVmO { currVm++; if (currVm surrVm=0; } allocatedVm(currVm): return currVm: }</pre>	<pre>if (currentAllocationCounts_size() < vmStatesList_size()){ for (int available/vmId : ymStatesList_keySet()){ if ('currentAllocationCounts_contains Key(available/vmId)){ ymId = available/vmId; break; } } } else { int currCount; int minCount = Integer MAX_VALUE; for (int this/vmId : currentAllocationCounts_keySet()){ currCount = currentAllocationCounts_get(this/ymId); if (currCount < minCount){ minCount = currCount; ymId = this/vmId; } allocatedVm(ymId); return ymId;</pre>	<pre>public int getNextAvalableVmO{ int ymld = -1; if (ymStatesList.size() > 0){ int temp; for (Iterator<integer> itr = ymStatesList.keySet0.iterator(); itr.hasNext();{ temp = itr.next(); VirtualMachineState state = ymStatesList.get(temp); //System.out.printh(temp + " state is " + state + " total yms " + ymStatesList.get(temp); if (state.equals(VirtualMachineSt ate.AVAILABLE)){ ymld = temp; break; } } allocatedVm(ymld); return ymld; }</integer></pre>	<pre>public int getNextAvailableVm(){ int ymld = -1; if(currentAllocationCounts.size() < ymStatesList.size()+1){ for (int availableVmld : ymStatesList.keySet()){ if (!currentAllocationCounts.contain: Key(availableVmld); ymld = availableVmld; break; } } } else { int currCount; int minCount = Integer MAX_VALUE; for (int thisVmld : currentAllocationCounts.keySet()){ { currentAllocationCounts.keySet(); { currentAllocationCounts.keySet(); { currentAllocationCounts.keySet(); { currentAllocationCounts.keySet(); { currentAllocationCounts.keySet(); { currentAllocationCounts.get(thisVmld;); if(currCount = currentAllocationCounts.get(thisVmld;); if(currCount < minCount)(minCount = currCount; ymld = thisVmld;); return ymld; } </pre>

In the above table, we can see the performance of existing algorithms for implementing the performance of response time and data transfer cost but we can analyze that performance of response time and data transfer cost not be in improving way at same time. Our proposed DSBP algorithm is basically implementation of throttled load balancing algorithm within cloud environment to achieve better response time and data transfer cost.

Proposed solution: Our proposed DSBP algorithm selects the data center according to job size, the expected processing time, network latency and availability of bandwidth to reduce the overall response time and processing time. The routing policy includes different parameters like, request size, user size, number of data centers, number of virtual machines, bandwidth, response time, data transfer cost, etc. The processing time depends upon the data center specification such as RAM, CPU and VM configuration. In Table 2, we have implemented the methods of round robin, active mentoring and throttled load balancing algorithm. Our proposed DSBP algorithm is the implementation of throttled load balancing algorithm. As a result, our proposed algorithm accommodates the current needs by taking real time values to evaluate the processing time to reduce the time needed to make the forwarding decision by the broker. The job processing time can depend on the computational task to be performed. For instance, a smaller job requires minimum processing time if there was no any I/O operation involved. However, since, it is not the service broker functionality to analyze the jobs and examine their complexity, we have considered the job size as an indication to the needed processing time with a positive relation between them. In Table 2, we present our implemented methods of existing and proposed DSBP algorithms. The proposed DSBP algorithm is the implementation of throttled load balancing algorithm, which works efficiently under the capacity of each virtual machine is different because the hardware configuration of virtual machines is different.

RESULTS AND DISCUSSION

Experimental setup and result analysis: In our experiment, multiple users are capable to use distributed data base center by using cloud based scheduling algorithms. For compute overloading, it is more useful to distribute load across the data centers over the parameters, bandwidth, VM image size, VM storage, VM memory, etc. for data center workload, a large number of users would be served by a cloud analyst where result

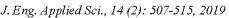
can be easily persisted on defined parameters. There are three principles presented in cloud analyst, User Base (UB), data center controller and internet. In cloud analyst, world is divided into 6 regions and the primary substances are user bases and data centers have a place to one of regions. This topographical collection is utilized to keep up a level of reasonable for the extensive scaled reproduction being endeavored in the cloud analyst. The cloud analyst internet presents reality of internet which highlights transmission dormancy and information exchange delay. Our simulation setup are carried out for 60 min by taking different number of users in form of 50 UB and 6 data centers from DC1-DC6. Each data center has 5 virtual machines separately with different parameters (Fig. 1-3).

The above configuration and simulation screens presents the performance of our proposed DSBP algorithm based on defined parameters in Table 3. This simulation result provides processing time, over all response time and data transfer cost in distributed data center environment experimented by cloud analyst. We have also taken this type of simulation results for existing round robin, active monitoring and throttled load balancing algorithms. After the performance all existing algorithms and our proposed DSBP, we get different type of values for processing time, response time and data transfer cost as given in Table 4, 5 and Fig. 4, 5.

			a content c	onfiguration					
Right	Simulation D	ur 60.0	min						
Simulation	User ba	Name	Region	Requests User	Data Size per Fequ (by.es)	Peak Hours Start (GHT)		Avg Peak Users	Avg OT-Fe Users
Graph		UPC	2			2		10000000	100
		UB2	2	60			9		100
		UB3	2	60				10000000	
		UB4	2	00			.9	100000	
Exit	Application	UB5		optimise ne				1000000	100
Esit	Application Deployment Configurati	Service Bro	kor Pell	optimise ne	sponse	-			
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Exit	Deployment	Service Brol Data Ce DC.	kor Pell	optimise ne	sponse	ge Size		y 512	BW 10
Exit	Deployment	Service Bro	kor Pell	optimise ne	sponse	ge Size		y	BW
Esit	Deployment	Data Ce	kor Pell	optimise ne	sponse	ge Size		512 512	BW 10 10

Fig. 1: Main configuration screen of userbase and application deployment

Main Configure Simulation Advanced Data Centers: Name Region Arch OS VMM Cost per VM SHr Memory Storage Data Physical Data Centers: Name Region Arch OS VMM Cost per VM SHr Memory Storage Data Physical DC1 0x86 Linux Xen 0.1 0.05 0.1 0.1 2 DC3 0x86 Linux Xen 0.1 0.05 0.1 0.1 1 DC4 0x86 Linux Xen 0.1 0.05 0.1 0.1 1 DC5 0x86 Linux Xen 0.1 0.05 0.1 0.1 1			J. 1	ing.	Аррие	cu ber	., 17 (.	2). 50		, 201	9	
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Centers: Name Region Addi OS VMM Cost per inemoty Storage Data Printal DC1 0x86 Linux Xen 0.1 0.05 0.1 0.1 22 DC2 0x86 Linux Xen 0.1 0.05 0.1 0.1 22 DC3 0x86 Linux Xen 0.1 0.05 0.1 0.1 1 DC3 0x86 Linux Xen 0.1 0.05 0.1 0.1 1 DC4 0x86 Linux Xen 0.1 0.05 0.1 0.1 1	Main Configurat	tion Data C	enter Confi	guration	Advanced	1						
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Centers: VM \$iHr Cost \$is Cost \$is Transfer Cost \$is HW Cost \$iGb Add New DC1 0x86 Linux Xen 0.1 0.05 0.1 0.1 22 DC2 0x86 Linux Xen 0.1 0.05 0.1 0.1 22 DC3 0x86 Linux Xen 0.1 0.05 0.1 0.1 10 DC3 0x86 Linux Xen 0.1 0.05 0.1 0.1 10 DC4 0x86 Linux Xen 0.1 0.05 0.1 0.1 0.05	Data	Name	Region	Arch	OS	VMM	Cost per	Memory	Storage	Data	Physical	
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DC4 0x86 Linux Xen 0.1 0.05 0.1 0.1 0		DC2	0	x86	Linux	Xen	0.1	0.05	0.1	0.1	1	Remove
		DC3	0	x86	Linux	Xen	0.1	0.05	0.1	0.1	1	
DC5 0x86 Linux Xen 0.1 0.05 0.1 0.1 1		DC4	0	x86	Linux	Xen	0.1	0.05	0.1	0.1	0	
		DC5	0	x86	Linux	Xen	0.1	0.05	0.1	0.1	1	



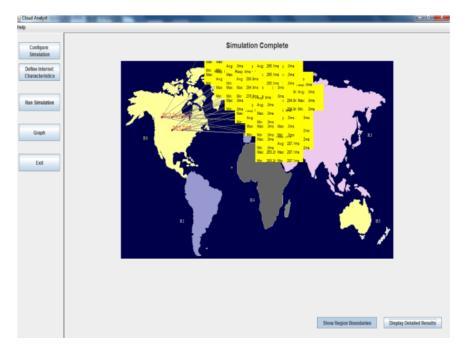
Data	Name	Region	Arch	OS	VMM C	Cost per	Memory	Storage	Data	Physical	
Centers:					١	VM \$/Hr	Cost \$/s	Cost \$/s	Transfer	HW	
									Cost \$/Gb	Units	Add New
	DC1	0	x86	Linux X	en	0.1	0.05	0.1	0.1	2	
	DC2	(x86	Linux X	en	0.1	0.05	0.1	0.1	1	Remove
	DC3		x86	Linux X	en	0.1	0.05	0.1		1	
	DC4		x86		en	0.1	0.05	0.1			2
	DC5	(x86	Linux X	en	0.1	0.05	0.1	0.1	1	
				Phy	sical Hardwa	are Details	s of Data Cer	nter : DC2			
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	Id		(Mb)	Storage (Mb)	Availab BW	ole I F	Number of	Proces	ed	Policy	

Fig. 2: Main configuration screen of data center

O C O			
Configure Simulatio	n		
Main Configuration Data Center Cor	figuration Advanced		
User grouping factor in User Ba (Equivalent to number of simult users from a single user base)	aneous		
Request grouping factor in Data (Equivalent to number of simult requests a single application so instance can support.)	aneous		
Executable instruction length p (bytes)	er request: 100		
Load balancing policy across VMPs in a single Data Ce	DSBP	¥	

Fig. 3: Configure simulation screen for proposed DSBP algorithm

Table 3: Parameters and values	
Parameters	Values
V-M image size	10000
VM memory	1024 MB
VM bandwidth	1000
Data Center (DC)	6
Virtual Machine (VM) in each data center	5
DC memory per machine	1000 MbPS
MBPS DC storage per machine	1GB
DC available bandwidth per machine	1000000



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Fig. 4: Simulation (running) screen of distributed data center

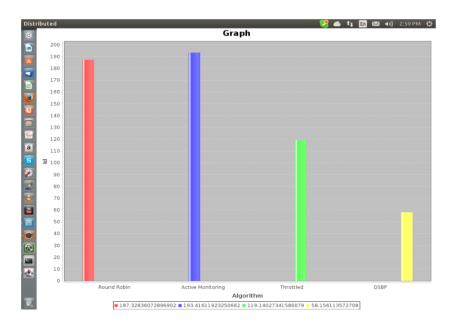


Fig. 5: Simulation graph of distributed data center for load balancing

Table 4: Comparison of overall response time and Data Center (DC) processing time summary
Distributed data center over all Response Time (RT)

	Round ro	obin active	(msec)	Active mor	nitoring algor	rithm (msec)	Throttle	d algorithm	(msec)	DSBP al	gorithm (n	1sec)
Algorithm/												
Response time	Avg.	Min.	Max.	Avg.	Min.	Max.	Avg.	Min.	Max.	Avg.	Min.	Max.
Over all RI	108.09	12.38	68.15	111.17	9.25	70.77	85.50	8.75	69.11	54.72	0.01	60.38
DC processing	98.3	0.00	0.22	72.19	0.00	0.22	60.01	43.23	0.22	0.12	0.03	0.20

Table 5: Compa Distributed data			4	(RT)								
	Round robin active (msec) Active monitoring algorithm (rithm (msec)	Throttlee	d algorithm	(msec)	DSBP al	gorithm (msec)		
Algorithm/ Response time	Avg.	 Min.	Max.	Avg.	Min.	Max.	Avg.	Min.	Max.	Avg.	Min.	Max.
Over all RI	108.09	12.38	68.15	111.17	9.25	70.77	85.50	8.75	69.11	54.72	0.01	60.38
DC processing time	98.3	0.00	0.22	72.19	0.00	0.22	60.01	43.23	0.22	0.12	0.03	0.20

CONCLUSION

For compute overloading, it is more useful to distribute across load of users request rather than within each data center over the different parameters (data center storage, memory, bandwidth, etc.) for data center workload. A large number of users would be served by a cloud analyst where result can be easily persisted on defined parameters. Our proposed solution is expected to have significant impact on the performance of distributed data center by virtue of its capabilities to adopt cloud based policies and load balancing algorithms. Load balancing and cloud collaboration drive the best solution in performance of distributed data center environment. Our proposed DSBP algorithm is able to play an important role to reflect the load balancing of different data centers for information processing at different level of logistics in real world computing environment. Hence, we conclude DSBP algorithm has effective result in logistics operations. We can also see an analytical performance of DSBP which has better results in comparison to round robin, active monitoring and throttled load balancing algorithms at a same time for reducing overall response time, data transfer cost and load in the distributed data center. The significance of this research is the demonstration of the DSBP algorithm using a cloud collaboration of advanced technologies which will surely help to present efficient performance of logistics management for tracking the information and managing the operations to distribute information from different location in distributed data center at same time. Future research is related to design a new load balancing algorithm for better resource utilization, minimum response time and minimum cost for fast throughput of cloud computing environment for logistics management to also maintain inventory distribution, vehicle locations etc.

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